On the Number of the Cusps of Rational Cuspidal Plane Curves

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Abstract

A cuspidal curve is a curve whose singularities are all cusps, i.e. unibranched singularities. The article describes computations which lead to the following conjecture:

A rational cuspidal plane curve of degree greater or equal to six has at most three cusps. The curves with precisely three cusps occur in three series.

Assuming the Flenner–Zaidenberg rigidity conjecture the above conjecture is verified up to degree 20.

Let \( C \subset \mathbb{P}^2 \) be a plane curve. The curve is called cuspidal if all its singularities are cusps, i.e., unibranched singularities. It is known that rational cuspidal plane curves underlie severe restrictions. For example Tono proved [T]:

**Theorem 1** A rational cuspidal plane curve has at most eight cusps.

It is not known whether this bound is sharp. In fact, one of the main purposes of this article is to provide evidence for the following conjecture:

**Conjecture 2** A rational cuspidal plane curve has at most three cusps — with the exception of a rational plane quintic with four cusps.

We will describe the topological type of a cusp singularity by its multiplicity sequence, \( \pi = (m_0, m_1, \ldots, m_n) \): Let

\[
V := X_n \xrightarrow{\pi_n} X_{n-1} \rightarrow \cdots \rightarrow X_1 \xrightarrow{\pi_1} X_0 = (\mathbb{C}^2, 0)
\]

\[
C_n \rightarrow C_{n-1} \rightarrow \cdots \rightarrow C_1 \rightarrow C_0 = (C, p)
\]

be the minimal embedded resolution of the curve germ \((C, p)\) such that the total inverse image \( D := \pi_n^{-1} \circ \cdots \circ \pi_1^{-1}(C) \) of the curve is a normal crossing divisor. Then the multiplicity sequence consists of the multiplicities \( m_i \) of the strict transforms \( C_i \) of the curve germs. For example a simple cusp has the multiplicity sequence \((2, 1, 1)\). To shorten the notation, sequences of the same multiplicity in the multiplicity sequence are indicated by indices, e.g., \((7, 2, 2, 1, 1)\) means \((7, 2, 2, 1, 1)\). (Here we follow the notation of [dJP] and [MS89], i.e., the number of 1s at the end of the multiplicity sequence equals the smallest number greater than one in the multiplicity sequence. [FZ96, FZ00] add an additional
1.) The ending sequence of 1s is usually dropped from the notation, because it can be recovered from the rest.

In this article we are hunting for cuspidal curves with the maximal number of cusps. The known rational cuspidal plane curves with at least three cusps are:

1. ([MS89, 4] or [Na]) For rational cuspidal plane curves of degree 5 the following combinations of cusps are possible: [(3), (2,2), (2)], [(2,2), (2,2), (2)], and [(2,3), (2), (2), (2)].

2. [FZ96, 3.5] For any \( d \geq 4, n \geq m \geq 1 \) with \( n+m = d-2 \) there is a rational cuspidal plane curve \( C \) of degree \( d \) with the three cusps \( [(d-2), (2n), (2m)] \).

3. [FZ00, 1.1] For any \( d = 2k + 3, k \geq 1 \), there is a rational cuspidal plane curve \( C \) of degree \( d \) with the three cusps \( [(d-3, 2k), (3k), (2)] \).

4. [F, 4] For any \( d = 3k + 4, k \geq 1 \), there is a rational cuspidal plane curve \( C \) of degree \( d \) with the three cusps \( [(d-4, 3k), (4k, 2), (2)] \).

All the above curves are unique up to projective equivalence. Not only is the existence of these curves known, but explicit constructions of them as well. Flenner, Zaidenberg, and Fenske described how to successively obtain the above curves in 2), 3), and 4) by plane Cremona transformations starting from three curves given by equations.

Flenner and Zaidenberg proved more: Assuming that a rational cuspidal plane curve \( C \) of degree \( d \geq 6 \) has at least three cusps and a point of multiplicity \( d-2 \) or \( d-3 \), then it is one of the curves in 2) or 3) above. Fenske showed the analogous statement with \( d-4 \) for the curves in 4) under the additional assumption of \( \gamma_2 = 0 \). This condition is explained below and for example implied by the (unproven) Flenner–Zaidenberg rigidity conjecture.

The purpose of this article is to explain how to obtain by computer computations the following Theorem:

**Theorem 3** If \( C \subseteq \mathbb{P}^2 \) is a rational cuspidal plane curve of degree \( d \leq 20 \) with at least three cusps and \( \gamma_2 = 0 \), then it is contained in the above list.

One is immediately lead to conjecture the following:

**Conjecture 4** Any rational cuspidal plane curve with at least three cusps is contained in the above list.

Finally, we explain the condition \( \gamma_2 = 0 \). Let \((V, D)\) be the minimal embedded resolution of \( C \subseteq \mathbb{P}^2 \), which was described above locally. Here \( D \) is a normal crossing divisor endowed with its reduced structure. The deformations of the pair \((V, D)\) correspond to equisingular deformations of the curve \( C \subseteq \mathbb{P}^2 \) [FZ96, 1.6–2.1]. Let \( \Theta_V \langle D \rangle \) denote the logarithmic tangent bundle. The infinitesimal automorphism group of the pair \((V, D)\) is \( H^0(V, \Theta_V \langle D \rangle) \), the infinitesimal deformations are \( H^1(V, \Theta_V \langle D \rangle) \), and the obstruction space for these deformations is \( H^2(V, \Theta_V \langle D \rangle) \).

We will assume throughout the article that the curve \( C \) has at least three cusps. Then \( X := \mathbb{P}^2 \setminus C \) is of log–general type [W]. By Iitaka's Theorem
[I, Theorem 6] the automorphism group of a log–general surface is finite, thus $H^0(V, \Theta_V(D)) = 0$. The pair $(V, D)$ or the curve $C$ is called rigid resp. unobstructed if $h^1(V, \Theta_V(D)) = 0$ resp. $h^2(V, \Theta_V(D)) = 0$. Flenner and Zaidenberg computed the Euler–characteristic of the logarithmic tangent bundle as

$$\chi(\Theta_V(D)) = -h^1(V, \Theta_V(D)) + h^2(V, \Theta_V(D)) = K_V(K_V + D),$$

where $K_V$ is the canonical divisor of $V$. They conjectured in an even more general setting that $(D, V)$ is always rigid and unobstructed. In our special situation Tono showed \[T, 4.1/5\]

$$\gamma_2 := h^0(2(K_V + D)) = K_V(K_V + D),$$

in particular $\gamma_2 = \chi(\Theta_V(D)) \geq 0$ and unobstructedness implies rigidity.

How to obtain the Theorem 3

The basic idea is to produce enough combinatorial restrictions on the multiplicity sequences of the cusps on a rational plane curve such that only the curves in the list remain as possible curves. These restrictions will be checked by a computer program [P].

The first step of the program is to produce all possible multiplicity sequences of cusps which might occur on a curve of degree $d$. By Proposition [FZ96, 1.2] a sequence of positive integers $m = (m_0, m_1, \ldots, m_n)$ is the multiplicity sequence of a cusp $p$ if and only if for each $i = 1, \ldots, n$ there is a $k \geq 0$ such that

$$m_{i-1} = m_i + \ldots + m_{i+k}, \quad \text{where} \quad m_i = m_{i+1} = \ldots = m_{i+k-1}$$

and the number of ending 1s equals the smallest $m_i > 1$.

The resulting curve has to be rational, hence it has to pass the following test:

- **$\delta$–Test**: The $\delta$–invariant of a cusp $p$ is $\delta_p = \sum_{i=0}^n m_i(m_i - 1)/2$. On a rational curve the sum of the $\delta$–invariants of the cusps equals the the arithmetical genus of $C$, which is $(d-1)(d-2)/2$. (See for example [OZ, Lemma 4].)

In particular, for any cusp we have $m_0 \leq d-1$. In fact, $m_0 \leq d-2$ because we want to have more than one singularity on $C$.

By Bezout’s Theorem the tangent line $T_pC$ intersects $C$ in $p$ with multiplicity at most $d$, i.e., $m_0 + m_1 \leq d$. However, if $m_0 + m_1 = d$ then $T_pC$ intersects $C$ only in $p$ and by the Lin–Zaidenberg Theorem [LZ] the affine curve $C \setminus \{p\} \subset \mathbb{P}^2 \setminus T_pC = \mathbb{C}^2$ is birational to a curve of the form $x^l - y^k = 0$ for some $k, l \in \mathbb{N}$ with $\gcd(l, k) = 1$. In this case $C$ has at most two cusps, thus in our case we must have $m_0 + m_1 < d$.

After producing the list of singularities which might occur on a rational cuspidal curve with at least three cusps, the program starts to search for possible combinations of singularities. First, it searches by a back–tracking algorithm for combinations of three up to eight cusps, such that the sum of their $\delta$–invariants equals the arithmetical genus; thereby implying that the resulting curve is rational. On these combinations the following tests are applied, which were already used by Flenner and Zaidenberg:
• \( \gamma_2 \)-Test: We examine only curves with \( \gamma_2 = 0 \). \( \gamma_2 \) can be computed as
\[
\gamma_2 = 3(3 - d) + \sum_{p \in \text{Sing} \, C} (\eta_p + \omega_p - 2),
\]
where \( \eta_p \) is defined for a cusp \( p \) with multiplicity sequence \( \overline{m} \) by
\[
\eta_p = \sum_{i=0}^{n} (m_i - 1)
\]
and \( \omega_p \) is the number of subdivisional blow-ups in the minimal resolution of the cusp \( p \) [FZ96, 2.5]. A blow-up is called subdivisional if the blow-up is performed at an intersection point of exceptional divisors of previous blow-ups; otherwise it is called sprouting. The exception is the first blow-up which is subdivisional by definition. (Note that we follow the definition of [MS89], where the \( \omega_p \) is one bigger than that of [FZ96, FZ00].) In terms of the multiplicity sequence the number \( \omega_p \) can be computed as ([FZ96, 2.3] or [MS89, Lemma 1])
\[
\omega_p = \sum_{i=1}^{n} \left( \left\lceil \frac{m_i - 1}{m_i} \right\rceil - 1 \right) + 1.
\]

• Bezout-Test: Each line in the plane intersects the curve in at most \( d \) points — counting multiplicities. In particular, for a line through the cusps \( p, q \), the first elements of their multiplicity sequences \( \overline{m}^p \) and \( \overline{m}^q \) must satisfy
\[
m_0^p + m_0^q \leq d.
\]

• Hurwitz-Test: Let \( p_1, \ldots, p_r \) denote the cusps of the curve and \( \overline{m}^{p_i} \) their multiplicity sequences. The projection from \( p_i \) is a \( d - m_0^{p_i} \) covering \( C \cong \mathbb{P}^1 \to \mathbb{P}^1 \), to which we apply the Riemann–Hurwitz formula. This yields
\[
2(d - m_0^{p_i}) - 2 = \sum_{q \in \overline{C}} (\nu_q - 1),
\]
where \( \nu_q \) is the ramification index of the covering in \( q \in C \). For the test we use the estimates \( \nu_{p_i} \geq m_1^{p_i} \) and \( \nu_{p_j} \geq m_0^{p_i} \) for \( j \neq i \), thus
\[
\sum_{j=1}^{r} (m_0^{p_j} - 1) + m_0^p + m_1^p \leq 2d - 2.
\]

These tests sum up the method of Flenner and Zaidenberg for the classification of rational cuspidal plane curves with at least three cusps and among them a cusp of multiplicity \( d - 2 \) or \( d - 3 \). They are also enough to prove our Theorem for \( d \in \{6, 8, 10, 14, 16\} \), i.e., the program described above produced only the known combinations of singularities in these cases. For the other degrees some additional combinations appeared. The easiest way to reduce the number of possible combinations further is to apply two propositions of Fenske. He observed that if the curve has a cusp whose multiplicity sequence starts like \([\lfloor d/2 \rfloor, \lfloor d/2 \rfloor, \lfloor d/2 \rfloor, \ldots] \) for odd \( d \) resp. \([d/2, d/2, d/2, \ldots] \) for even \( d \), then it can be transformed into another rational cuspidal plane curve by a plane Cremona transformation. More precisely, his results are:
1. [F, 4.1.13, 4.1.14] A rational cuspidal curve $C$ of degree $d = 2k + 1$ and cusp type $[(k, k, k, m_1), m_2, \ldots, m_r]$ can be taken by a Cremona transformation to a rational cuspidal curve $C'$ of degree $d' = k + 2$ and cusp type $[m_1, m_2, \ldots, m_r]$. Further, the curve $C'$ has a regular point $q$ whose tangent line cuts $C'$ with multiplicity 2 in $q$ and with multiplicity $k$ in the cusp of type $m_1$.

Rational cuspidal curves of degree $d = 2k + 1$ and cusp type $[(k, k, k, m_1), m_2, \ldots, m_r]$ with $r \geq 3$ and $m_2, \ldots, m_r$ non-trivial do not exist.

2. [F, 4.1.15] A rational cuspidal curve $C$ of degree $d = 2k$ and cusp type $[(k, k, k, m_1), m_2, \ldots, m_r]$ can be taken by a Cremona transformation to a rational cuspidal curve $C'$ of degree $d' = k$ and cusp type $[m_1, m_2, \ldots, m_r]$. Further, the tangent line to $C'$ in the cusp of type $m_1$ intersects $C'$ only in this cusp (with multiplicity $k$).

This excludes the following number of cases:

<table>
<thead>
<tr>
<th>$d$</th>
<th>number of cases</th>
</tr>
</thead>
<tbody>
<tr>
<td>7</td>
<td>3</td>
</tr>
<tr>
<td>9</td>
<td>5</td>
</tr>
<tr>
<td>11</td>
<td>9</td>
</tr>
<tr>
<td>12</td>
<td>2</td>
</tr>
<tr>
<td>13</td>
<td>16</td>
</tr>
<tr>
<td>15</td>
<td>6</td>
</tr>
<tr>
<td>17</td>
<td>52</td>
</tr>
</tbody>
</table>

This deals with most previously unknown combinations. The remaining ones are listed in the following table, together with the name of a test, which shows that it cannot exist. We will describe these tests below.

<table>
<thead>
<tr>
<th>$d$</th>
<th>cusp type</th>
<th>Test</th>
</tr>
</thead>
<tbody>
<tr>
<td>13</td>
<td>$(8, 4_0), (2), (2)$</td>
<td>$\Gamma$–Test with $l = 3$ for $(8, 4_0)$</td>
</tr>
<tr>
<td>17</td>
<td>$(7_5, 5, 2_2), X_3$</td>
<td>$\Gamma$–Test with $l = 2$ for $(7_5, 5, 2_2)$</td>
</tr>
<tr>
<td>18</td>
<td>$(8_4, 6, 2_7), (2), (2)$</td>
<td>$\Gamma$–Test with $l = 2$ for $(8_4, 6, 2_7)$</td>
</tr>
<tr>
<td>18</td>
<td>$(8_4, 6, 2_6), X_3$</td>
<td>$\Gamma$–Test with $l = 2$ for $(8_4, 6, 2_6)$</td>
</tr>
<tr>
<td>18</td>
<td>$(8_4, 6, 2_5), X_4$</td>
<td>$\Gamma$–Test with $l = 2$ for $(8_4, 6, 2_5)$</td>
</tr>
<tr>
<td>18</td>
<td>$(8_4, 6, 2_4), \ldots$</td>
<td>$\Gamma$–Test with $l = 2$ for $(8_4, 6, 2_4)$</td>
</tr>
<tr>
<td>18</td>
<td>$(8_4, 6, 2_3), \ldots$</td>
<td>$\Gamma$–Test with $l = 2$ for $(8_4, 6, 2_3)$</td>
</tr>
<tr>
<td>18</td>
<td>$(8_4, 4_3, 3), X_3$</td>
<td>Double Covering / Spectrum–Test</td>
</tr>
<tr>
<td>18</td>
<td>$(8_4, 4_3, 2_4), (2), (2)$</td>
<td>Double Covering / Spectrum–Test</td>
</tr>
<tr>
<td>18</td>
<td>$(8_4, 4_3, 2_3), X_3$</td>
<td>Double Covering / Spectrum–Test</td>
</tr>
<tr>
<td>18</td>
<td>$(8_4, 4_3, 2_2), X_4$</td>
<td>Double Covering / Spectrum–Test</td>
</tr>
<tr>
<td>19</td>
<td>$(12_6_5, 3_2, 3), (2), (2)$</td>
<td>Spectrum–Test</td>
</tr>
<tr>
<td>19</td>
<td>$(12_6_5, 3_3), X_3$</td>
<td>Spectrum–Test</td>
</tr>
<tr>
<td>19</td>
<td>$(12_6_5, 4_2), (2), (2)$</td>
<td>$\Gamma$–Test with $l = 3$ for $(12_6_5, 4_2)$</td>
</tr>
<tr>
<td>19</td>
<td>$(12_6_5, 4_2), X_3$</td>
<td>$\Gamma$–Test with $l = 3$ for $(12_6_5, 4_2)$</td>
</tr>
<tr>
<td>19</td>
<td>$(12_6_5, 4_2), X_4$</td>
<td>$\Gamma$–Test with $l = 3$ for $(12_6_5, 4_2)$</td>
</tr>
<tr>
<td>20</td>
<td>$(9_1, 5_4), (3), (2)$</td>
<td>$\Gamma$–Test with $l = 2$ for $(9_1, 5_4)$</td>
</tr>
<tr>
<td>20</td>
<td>$(8_6), X_3$</td>
<td>Spectrum–Test</td>
</tr>
</tbody>
</table>

Here, $X_3$ and $X_4$ stand for one of the following combinations of cusps singularities:

$X_3 \in \{[(2_2), (2)], [(2), (2), (2)]\}$

$X_4 \in \{[(3), (2)], [(2_3), (2)], [(2_2), (2)], [(2_2), (2), (2)], [(2), (2), (2), (2)]\}$

The proof of the following simple test is based on Bezout’s Theorem:
Proposition 5 (Γ-Test [FLMN, 3.4 Prop. 2]) Let \( \Gamma_i \) be the semi-group rings of the \( r \) cusps \( p_i \) on a cuspidal plane curve of degree \( d \). Then for all \( l < d \) and \( n_i \in \mathbb{N}_0 \) with \( \sum n_i \geq ld \) the following inequalities hold

\[
\sum_{i=1}^{r} \#(\Gamma_i \cap [0, n_i]) \geq (l + 1)(l + 2)/2.
\]

For the combinations in the table where this method is indicated it was enough to apply this proposition to the worst singularity alone to see that it cannot occur on a curve of that particular degree.

Another known method to exclude some combinations of singularities is through the cyclic covering method ([MS89, Sec. 3] or [Y1, Y2]). It works as follows: Assume that a plane curve \( C \) with a certain combination of singularities exists. One constructs a surface \( S \) as the minimal resolution of a cyclic covering of \( \mathbb{P}^2 \) branched along the curve \( C \) and then shows that this surface has impossible invariants. Even using the algorithm of Némethi to construct the resolution of the covering [Ne, §5] this is a very laborious method and cannot deal with all the remaining cases.

The most powerful method seems to be provided by the semi-continuity of the spectrum ([FLMN, 3.6] or [K, 8.9]). If \((C, p_i)\) are the singularities of a plane curve \( C \) of degree \( d \), then the multispectrum \( \sum_i (C, p_i) \) is a deformation of the plane germ \((x^d + y^d, 0)\). Therefore, the union of the spectra of the singularities \((C, p_i)\) satisfies the semi-continuity property compared with the spectrum of \((x^d + y^d, 0)\) for any interval \((\alpha, \alpha + 1]\). The spectrum of a plane cusp can be computed by purely combinatorial means using Saito’s formula [S]. The author implemented this formula into Singular [GPS] as part of the library \texttt{cusp.lib} which extends the built-in library \texttt{gmssing.lib} [P]. With the help of this library such checks can be easily done.

References


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