

The number of non-solutions to an equation in a group and non-topologizable torsion-free groups

(Klyachko and Trofimov, 2004)

① Definitions and Markov's theorem

Def. 1 A group G is called topologizable if it admits a non-discrete Hausdorff group topology.

Def. 2 A subset $S \subseteq G$ is called elementary algebraic if there exist elements $a_1, \dots, a_m \in G$ and numbers $\epsilon_1, \dots, \epsilon_n \in \mathbb{Z}$, s.t. $S = \text{Sol}(a_1 x^{\epsilon_1} a_2 x^{\epsilon_2} \dots a_m x^{\epsilon_m} = 1)$.

Thm (Markov, 1946) A countable group G is non-topologizable
 $\Leftrightarrow G \setminus \{1\}$ is a finite union of elementary algebraic sets

$\Leftrightarrow \exists$ a finite subset $F \subseteq G$:

$G \setminus F$ is a finite union of elementary algebraic sets

Dem Let X be a Hausdorff top. space and let
 $F = \{x_1, \dots, x_n\} \subseteq X$ be a finite subset.

Then F is open \Leftrightarrow each $\{x_i\}$, $i=1, \dots, n$ is open.

Proof $\exists O(x_1) : x_2, \dots, x_n \notin O(x_1)$. Then

$$\{x_1\} = F \cap O(x_1).$$

Main theorem [KT]. \exists f.g. torsion-free group H

and one equation $w(x)=1$ with one variable and
coefficients from H , s.t. $H \setminus \{1\} = \text{Sol}(w(x)=1)$.

Remark

- 1) By Löwenheim-Skolem, f.g. can be replaced by any infinite cardinality.
- 2) We will use later that $\exp_x(w) = 0$.

Earlier was known

$$\text{Trofimov (2004): } G \setminus \{g_1, \dots, g_{2n}\} = \text{Sol}([x, a]^n = 1)$$

↓
 f.g., $\pi(G) = N \cup \{\infty\}$.

↑
 some element of G .

Thm 1 For any two cardinals s and n with $s+n=\infty$, there exists a group G of cardinality $s+n$ and an equation $u(x)=1$ over G , s.t.

$$|Sol(u(x)=1)| = s \quad \& \quad |Sol(u(x)\neq 1)| = n.$$

Exercise $(2,1)$ is not realizable: in \mathbb{Z}_3 any equation $u(x)=1$ has exactly 0, 1, or 3 solutions

Proof: Consider an equation $x^n = a$, where $\langle a \rangle = \mathbb{Z}_3$.
 If $x_1^n = a = x_2^n$, $x_1 \neq x_2$, then $(x_1 x_2^{-1})^n = e \Rightarrow n \mid 3 \Rightarrow n=1$
 $\Rightarrow 3$ solutions.

Thm 1 For any two cardinals s and n with $s+n=\infty$, there exists a group G of cardinality $s+n$ and an equation $u(x)=1$ over G , s.t.

$$|\text{Sol}(u(x)=1)| = s \quad \& \quad |\text{Sol}(u(x) \neq 1)| = n.$$

Proof 1) $n=0$. Equation $1=1$

2) $s=0$. $g=1$, where $g \in G \setminus \{1\}$.

3) $s=1$ $x=1$.

4) $1 < s \leq n \quad (\Rightarrow n \text{ is infinite})$

$$G = A * B$$

Take $a \in A \setminus \{1\}$. Then

$$\text{Sol}_G(xa=ax) = A$$

$$\text{Sol}_G(xa \neq ax) = G \setminus A.$$

5) $s > n > 0 \quad (\Rightarrow s \text{ is infinite})$

$$G = H \times K$$

08 Solutions : $(H \setminus \{1\}) \times K$
of $u(x)=1$

04 Nonsolutions : $\{1\} \times K$.

2021 Proof: $[a, hk] = [a, h]$

$\Rightarrow k$ does not influence.

From the proof of the main thm:

$$H = \langle a, b, c_1, \dots, c_e \mid R \rangle$$

Equation:

$$\underbrace{[c_1 v([a, xc]) c_1^{-1}, v([b, x])] = 1}_{\text{" } u(x) \text{"}}$$

where c_1, a, b are as above,
 v is a word, depending
on generators of H and one
variable y .

② Ol'shanskii's approach to constructing groups with prescribed properties

A graded presentation $G(\infty) = \langle A \mid R \rangle$ with a filtration on the set of relators : $R = \bigcup_{i=0}^{\infty} R_i$, $\emptyset = R_0 \subseteq R_1 \subseteq \dots$

Relators from $R_i \setminus R_{i-1}$ are called relators of rank i ,

$G(i) = \langle A \mid R_i \rangle$ is also considered as a graded presentation with $R_i = R_{i+1} = \dots$

Def. A graded presentation $G(\infty)$ satisfies condition R with parameters α, h, d, n if for each $i = 0, 1, 2, \dots$, there exists a set $X_i \subseteq F(A)$ of words (called periods of rank i), s.t.

1) $W \in X_i \Rightarrow |W| = i$, $W \not\sim (W')^m$ \leftarrow any $\in \mathbb{Z}$
 $G(i-1) \leftarrow$ any in $F(A)$ with $|W'| \leq i-1$.

2) $U \neq V$ from $X_i \Rightarrow U \not\sim V^{\pm 1}$
 $G(i-1)$

3) Each relator $R \in R_i \setminus R_{i-1}$ has the form

$$R = \prod_{k=1}^r (T_k A^{n_k}), \text{ where } A \in X_i \text{ and}$$

n_k, T_k, A satisfy :

R1 $|n_k| \geq n$.

R2 $|n_s|/|n_t| \leq 1 + \frac{1}{2}h^{-1}$. (think that $|n_s| = |n_t|$)

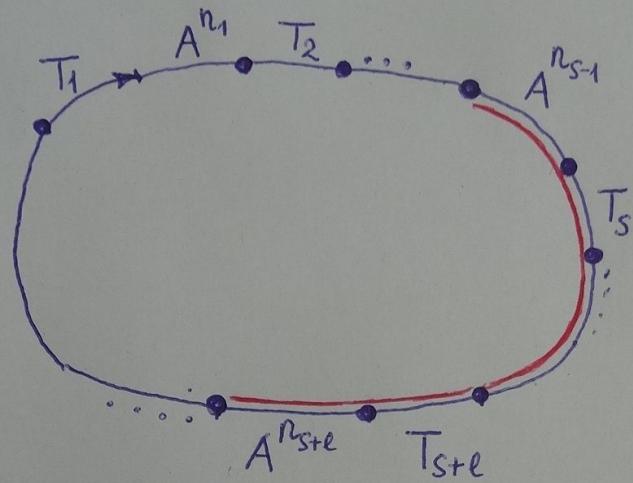
R3 The words T_k are not equal in $G(i-1)$ to words of smaller length, and $|T_k| \leq d_i$

R4 T_k is not a power of A_k in $G(i-1)$.

R5-R6

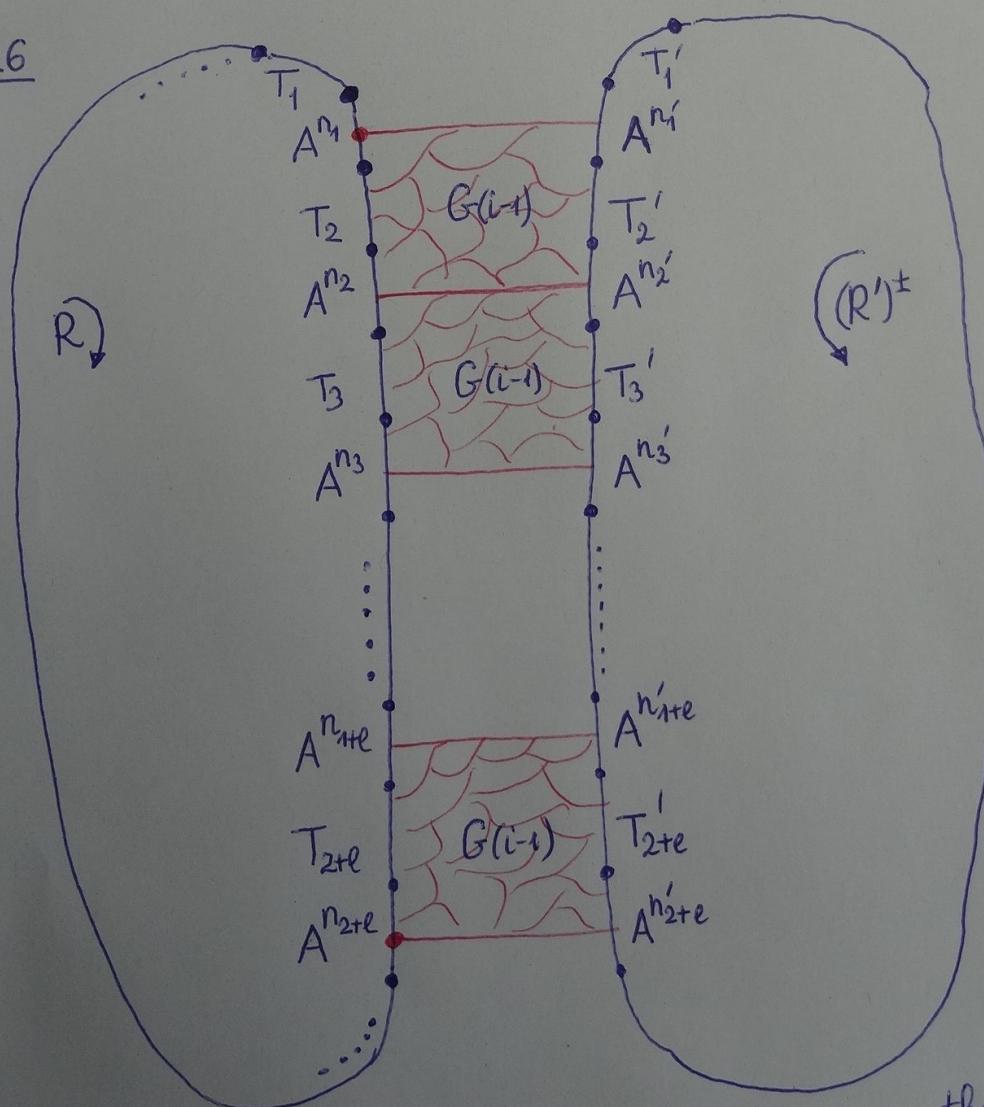
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 if $n_k \ll n$,
 then T_k is
 small and
 A^{n_k} large

R5. The word R is not a proper power in the free group.
and any long subword of ~~R~~^l the cyclic word R
occurs in R only once:



$l > \bar{\alpha}^1 - 4$
pairs of syllables

R6



Suppose:
 1) $l \geq \bar{\alpha}^1 - 2$,
 2) $\text{Sign}(n_i) = \text{Sign}(n'_i)$
 $i = 1, \dots, 2 + l$,
 3) A contiguity condition
 (see red)

Then $R \equiv R'$,

$V \equiv V'$



the subword
of R from one
red point to
the other

and V is not
a subword of
the cyclic word R^{-1}

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6)

Lem 1 If a graded presentation $G(\infty)$ satisfies cond. R
 (from Ol'sh.) with $1 \ll \alpha' \ll h \ll d \ll n$, then

Thm 26.5

1) all abelian subgroups of $G(\infty)$ are cyclic

Lem 25.2

2) the group $G(\infty)$ is torsion-free (in particular $|G(\infty)| = \infty$)

Lem 25.4

3) if $X \sim Y$ (from $G(\infty)$), then $\exists Z \in G(\infty): X = ZYZ^{-1}$,
 $|Z| \leq (\frac{1}{2} + \alpha) (|X| + |Y|)$

Lem 22.2

4) if A, B, C are nontrivial elements from $G(\infty)$, $X \in G(\infty)$
 such that $X^t A X B \sim C$, then the double coset
 $\langle A \rangle X \langle B \rangle$ contains an element X' of length
 $|X'| < (\frac{1}{2} + \alpha) (|A| + |B| + |C|) + [\frac{1}{2}|A|] + [\frac{1}{2}|B|]$

Lem 23.16

5) if a word $X = 1$ (from $G(\infty)$), then $X = 1$ in the group

$$\langle A \mid \{R \in \mathcal{R} \mid |R| < (1 - \alpha)^n \cdot |X|\} \rangle$$

Lem 25.14

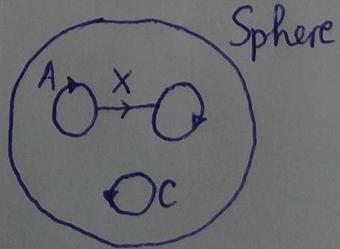
6) if $[X, ZXZ^{-1}] = 1$ (from $G(\infty)$), then $[X, Z] = 1$ (from $G(\infty)$)

Geometry for 4):

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Useful to read from the book of Ol'shanski:
 i) Pages 218-226 (B-maps) from Chapter 7.

ii) §25 from Chapter 8

1)

③ Construction of H

We fix $n \gg d \gg h \gg 1$
 ↪ even number.

Generators $A = \{a, b, c_1, \dots, c_h\}$.

We set $R_0 = R_1 = R_2 = \emptyset$ and define $G(i) = \langle A | R_i \rangle$, $i \geq 2$
 assuming the presentation $G(i-1) = \langle A | R_{i-1} \rangle$ is defined.

- Definition of periods $X_i \subseteq F(A)$.

This is a maximal subset satisfying:

$$1) W \in X_i \Rightarrow |W| = i$$

$$2) W \not\sim W' \text{ if } W \in X_i \text{ and } |W'| < |W| \\ G(i-1)$$

$$3) W_1 \neq W_2 \in X_i \Rightarrow W_1 \not\sim W_2^{\pm 1} \\ G(i-1)$$

$$4) W \in X_i \Rightarrow W = ab \cdot (c_1 c_2 \dots c_h)^l \text{ in } G(i-1) / [G(i-1), G(i-1)] \\ \text{for some } l \in \mathbb{Z}.$$

- Definition of $R_i = R_{i-1} \cup \widetilde{T}_i$:

\widetilde{T}_i is a maximal set of pairwise nonconjugate in $G(i-1)$
 words from $F(A)$ of the form

$$R = \prod_{j=1}^h (T_j A^{(\leftarrow 1)^{f_n}}), \text{ where } A \in X_i \text{ and}$$

T_1, \dots, T_h are words from $F(A)$ s.t. $\forall j=1, \dots, h$:

1) $|T_j| < d_i$ (small)

2) $T_j \neq u$ if $|u| < |T_j|$ (minimal)
 $G(i-1)$

3) $\exists Z \in F(A) : (T_1, \dots, T_h) \sim_{G(i-1)}^Z (c_1, \dots, c_h) Z$

Lem. 2 $H = G(\infty) = \langle A | R \rangle$, where $R = \bigcup_{i=0}^{\infty} R_i$ satisfies R .

Proof Mod $[F(A), F(A)]$ we have

- $R = c_1 c_2 \dots c_h$
- $T_j = c_j, j=1, \dots, h$
- $A = a \underset{\cap}{\underset{x_i}{\underset{\uparrow}{\underset{\wedge}{}}} b (c_1 c_2 \dots c_h)^l$ for some l
- $G(i)/[G(i), G(i)] = \langle a, b, c_1, \dots, c_h \mid c_1 c_2 \dots c_h = 1 \rangle, i \geq 2$
 $= G(\infty)/[G(\infty), G(\infty)]$

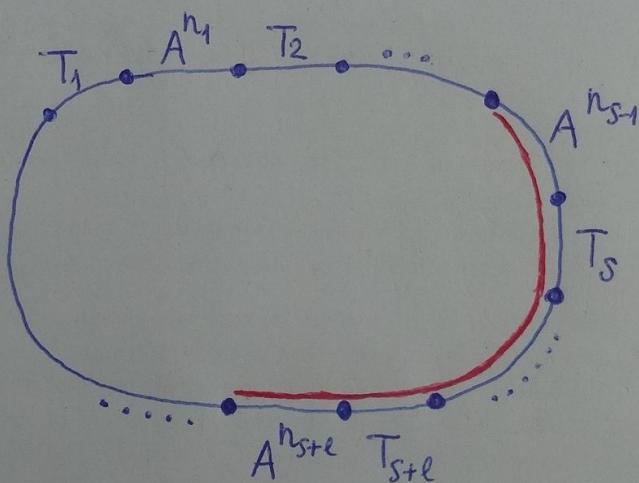
\Rightarrow $\underset{\uparrow}{\underset{x_i}{\underset{\wedge}{\underset{\cap}{}}} A$ is not a proper power mod ~~$G(i-1), [G(i-1), G(i-1)]$~~ in $G(i-1)$ and even in $G(i-1)/[G(i-1), G(i-1)]$

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\Rightarrow Properties 1), 2), $R_1 - R_2$ ~~from the definition of condition R~~ are satisfied.

Verification of R5.

R5 The word R is not a proper power in the free group and any long subword of the cyclic word R occurs in R only once:



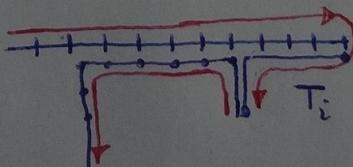
$l > \bar{\lambda}^{-1} - 4$
pairs of syllables

Attention: There may be cancellations in

$$R = T_1 A^{n_1} T_2 A^{n_2} \dots T_h A^{n_h}$$

But these cancellations are small:

In each triple $A^{n_{i+1}} T_i A^{n_i}$ if T_i cancels completely, then only a couple (or few) A's cancel further.

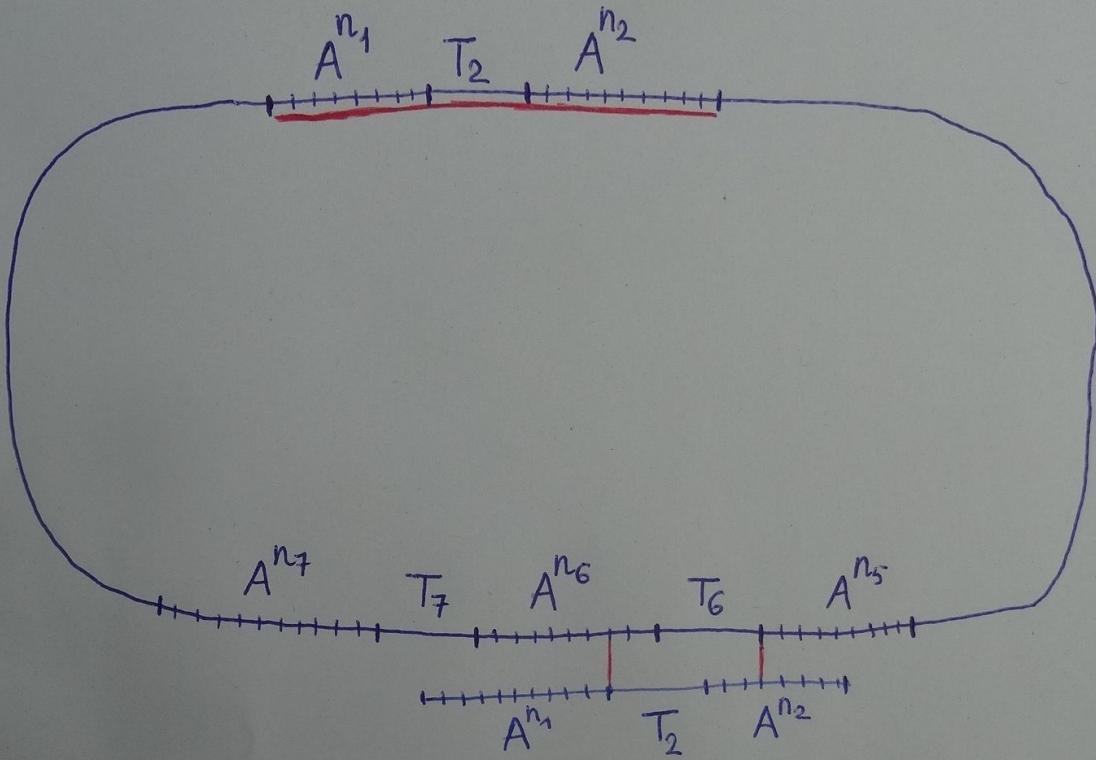


Use: long overlapping A-words
(powers of A)
must match
be consistent.



$T_i \in \langle A \rangle \leftarrow$ forbidden by R4.

Continue to verify R5



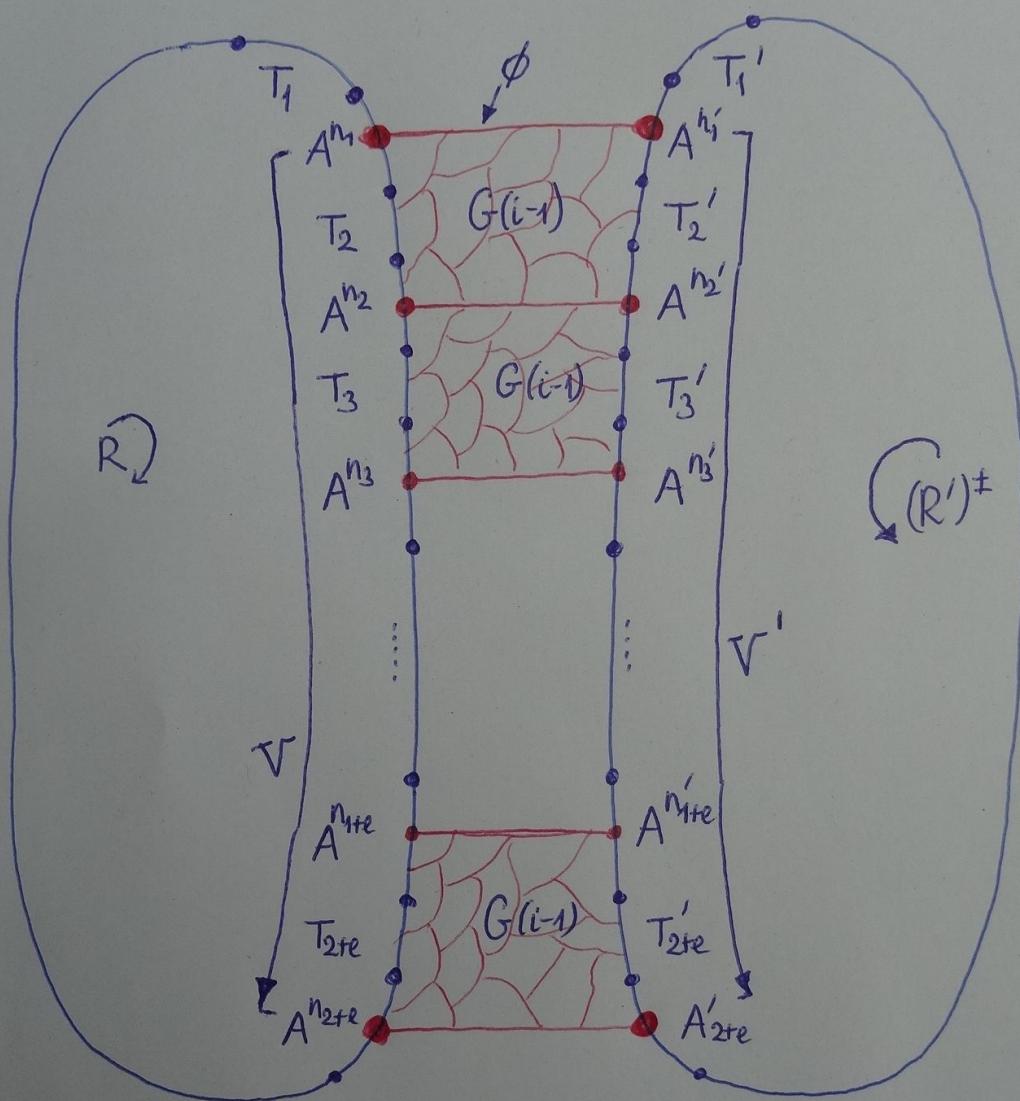
Because of matching, we have $A^s T_2 = T_6 A^t$ for some s, t .



Consider this equality mod $[F(A), F(A)]$:

$$(ab(c_1 c_2 \dots c_h)^e)^s \cdot c_2 = c_6 \cdot (ab(c_1 c_2 \dots c_h)^e)^t \quad \leftarrow \text{impossible.}$$

Verification of R6



Formulation

of R6: Suppose $l \geq \alpha^{-1} - 2$, $\text{sign}(n_i) = \text{sign}(n'_i)$, $i = 1, \dots, 2+l$ and there exist a subdivision ~~by red~~ of V and V' by red points such that the vertical sides of each of $l+1$ rectangles are equal in ~~red~~ $G(i-1)$

Then $R \equiv R'$, $V \equiv V'$ and V is not a subword of the cyclic word R^{-1}

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6)

We have equations of kind $A^\alpha T_j A^\beta \underset{G(i-1)}{=} A^{\alpha'} T_{j'}^{\pm 1} A^{\beta'} \quad (0)$

Rewrite: $T_j A^P \underset{G(i-1)}{=} A^S T_{j'}^{\pm 1} \quad (1)$

Consider mod $[G(i-1), G(i-1)] \Rightarrow j=j', p=s, \pm 1=1$

Recall that by def. of $T_j, T_{j'}$, there exist $Z, Z' \in F(A)$:

$$T_j = \underset{G(i-1)}{Z c_j Z^{-1}}, \quad T_{j'} = \underset{G(i-1)}{Z' c_j (Z')^{-1}}. \quad (*)$$

Substitute in (1):

$$\underset{G(i-1)}{Z c_j Z^{-1} A^P} = A^P \underset{G(i-1)}{Z' c_j (Z')^{-1}}$$

Rewrite: $c_j \underset{G(i-1)}{\leftrightarrow} \underset{G(i-1)}{Z^{-1} A^P Z'}$ (2)

By induction, $G(i-1)$ satisfies condition R, in particular commuting elements of $G(i-1)$ must lie in the same cyclic subgroup (Lem. 1). But c_j is not a proper power in $G(i-1)$ (again mod $[G(i-1), G(i-1)]$).

$$\Rightarrow \underset{G(i-1)}{Z^{-1} A^P Z'} = c_j^k \quad \text{for some } k \quad (3)$$

Consider $j+1, j+2, \dots$ instead of j and analogously (by ass. of R) we have many equalities of kind (0).

$$\underset{G(i-1)}{Z^{-1} A^{P_1} Z'} = c_{j_1}^{k_1}. \quad (4)$$

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$$\underset{G(i-1)}{Z^{-1} A^{P-P_1} Z} \underset{G(i-1)}{=} c_j^k c_{j_1}^{-k_1}. \quad \text{Again mod } [G(i-1), G(i-1)] \quad (5)$$

gives $p=p_1, k=k_1=0$. Then (3) implies $\underset{G(i-1)}{Z'} = A^P Z$. Then, by (*)

$$T_{j'}' = \underset{G(i-1)}{A^P T_j A^P} \Rightarrow R \underset{G(i-1)}{\sim} R'. \Rightarrow R = R' \quad (\text{see def } T_i)$$